

ETMF 2023 – FORMAL METHODS: WHAT ARE THEY? WHAT ARE THEY USED FOR?

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Lecturer @ CIn-UFPE, Brazil

Research interests: **formal methods** (theory and practice)



CIn-TRUST

trust.cin.ufpe.br
cin.ufpe.br/~ghpc



Project: **RoboTIC@**

ines.org.br/projects-2



ROBOSTAR

robostar.cs.york.ac.uk

RoboWorld: more information
(publications, examples, tool, etc.)

[robostar.cs.york.ac.uk/
roboworld](http://robostar.cs.york.ac.uk/roboworld)



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A connected (software-based) society



Software and systems should behave as expected

But, do we trust them?

Formal methods to the rescue!

Small digression: aeroplanes

Source: <https://ccntservice.airbus.com/apps/cockpits/a380>



© AIRBUS S.A.S. 2014 - photo by master films



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Small digression: aeroplanes

KLM cockpit tales: <https://www.youtube.com/watch?v=XXM01nnlnxs>



Small digression: aeroplanes

Flight tests: https://www.youtube.com/playlist?list=PLjl-u2y72YNwiW9CQ1Tp6yOsR-WkL_QIX



Formal methods

What about the **rigour** in:
software system **specification**, **design**, and **verification**?



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Formal methods

What about the **rigour** in:
software system **specification**, **design**, and **verification**?

*“Formal methods are **mathematical approaches** to software and system development which **support the rigorous specification, design and verification** of computer systems.”*

Formal Methods Europe – **FME**

An important role on achieving **high trustworthiness levels**



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Formal methods: not only about aeroplanes

AWS: How Amazon web services uses formal methods

<https://doi.org/10.1145/2699417>

Google: Using Formal Methods at Google

<https://datatracker.ietf.org/meeting/118/session/ufmrg>

Meta: Moving Fast with Software Verification

https://doi.org/10.1007/978-3-319-17524-9_1

Microsoft: SMT in Verification, Modeling, and Testing at Microsoft

https://doi.org/10.1007/978-3-642-39611-3_3

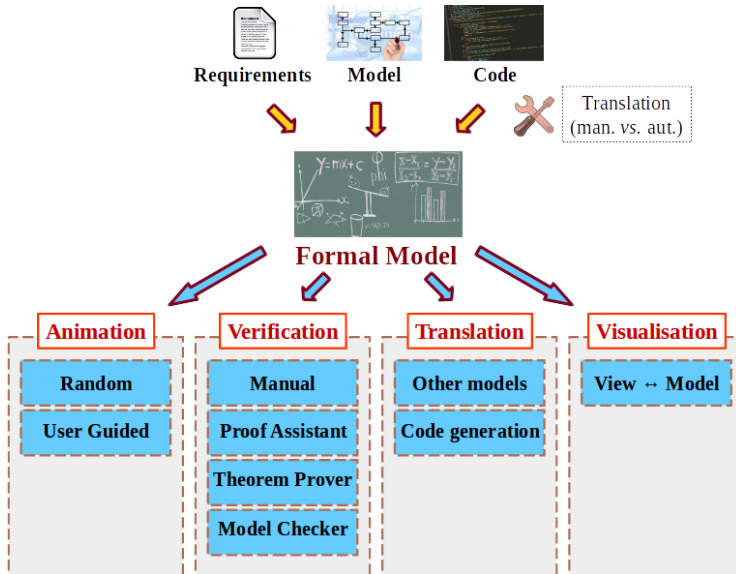


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Formal methods: a big picture



Agenda

1 B language

2 CSP

3 Timed automata

4 SMT-LIB v2

5 Gallina

6 Conclusion



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B language

Focus: safety-critical software (single threaded)

Application: numerous applications, particularly in the **railway domain**
(https://link.springer.com/chapter/10.1007/978-3-030-58298-2_8)

Tools

- **Atelier B:** <https://atelierb.eu/>
- **ProB:** <https://prob.hhu.de/w/index.php>
- **BMotionWeb:**
<https://prob.hhu.de/w/index.php?title=BMotionWeb>



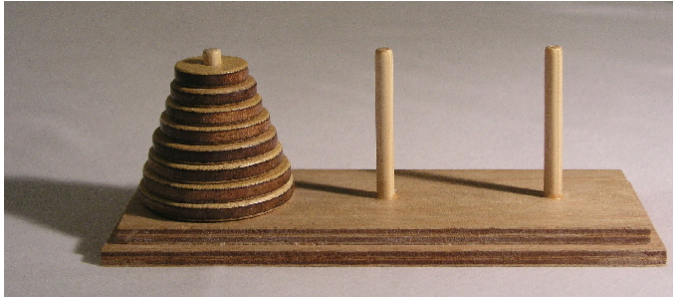
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B language: example

Tower of Hanoi puzzle¹



¹Source: [Wikimedia Commons](#)



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B language: specification

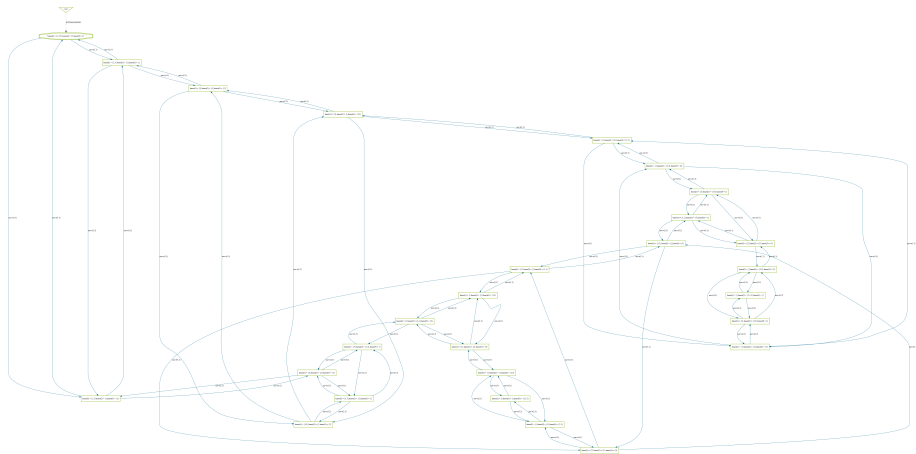
```

7- VARIABLES
8   hanoi
9- INVARIANT
10  // System state is described as a function from pegs (1..3) to disks (1..3)
11  // 1 denotes the smaller disk, whereas 3 denotes the biggest one
12  hanoi : 1..3 --> POW(1..3)
13  // There is no disk repetition between pegs
14  & ! (p1,p2) . ((p1 : dom(hanoi) & p2 : dom(hanoi) & p1 /= p2 => hanoi(p1) /\ hanoi(p2) = {}))
15  // Is there a solution?
16  & not((hanoi(1) = {}) & (hanoi(2) = {}) & (hanoi(3) = 1..3))
17- INITIALISATION
18  // Initially, all disks on the first peg
19  hanoi := {1 |-> {1,2,3}, 2 |-> {}, 3 |-> {}}
20- OPERATIONS
21  // Moving disks between pegs
22  move(from,to) =
23-  PRE
24    // The source and target pegs shall be different
25    from : dom(hanoi) & to : dom(hanoi) & from /= to
26    // The source peg is not empty
27    & hanoi(from) /= {}
28    // You cannot place a bigger disk over a smaller one
29    & (hanoi(to) /= {} => min(hanoi(from)) < min(hanoi(to)))
30-  THEN
31    hanoi := hanoi <+
32    // Updating the state of the source and target pegs
33-    {
34      from |-> (hanoi(from) - {min(hanoi(from))}),
35      to |-> (hanoi(to) \/ {min(hanoi(from))})
36    }
37-  END

```



B language: state space



28 states and 79 transitions
Graph generated with ProB

```
$ dot -Tpng Hanoi.dot > Hanoi.png
```



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B language: visualisation

Another example: `lift`²

Simple Lift System

Edit ProB Diagram View Window

Simulator Editor: visualization.svg

User Interactions Log

Transition	Executor
INITIALISATION()	
SETUP_CONSTANTS()	

Events

Filter Events

- move_up()
- move_down()
- door_open()
- door_close()
- ▶ switch_move_up()
- ▶ switch_move_down()
- ▶ switch_move_stop()
- ▶ send_request(-1)
- ▶ send_request(0)
- ▶ send_request(1)

²Source: [BMotionWeb examples](#)

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CSP

Focus: concurrent systems

Application: numerous applications

(https://doi.org/10.1007/978-3-319-51046-0_4)

Tools

- **FDR:** <https://cocotec.io/fdr/>



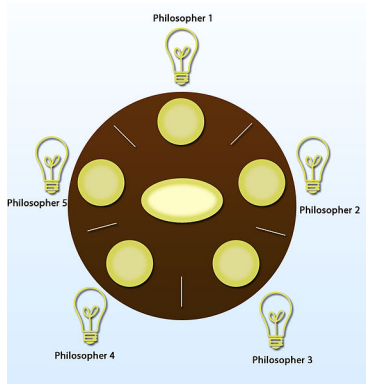
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CSP: example

Dining philosophers problem³



³Source: [Wikimedia Commons](#)

CSP: specification⁴

```

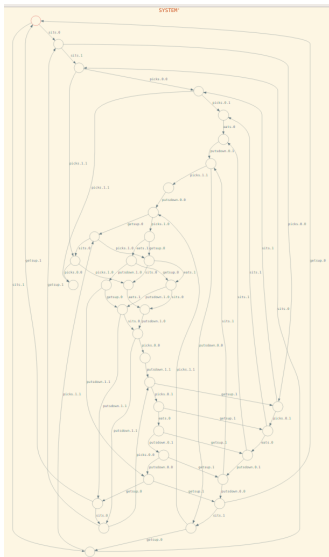
1  -- Bill Roscoe
2  N = 4
3
4  PHILNAMES= {0..N-1}
5  FORKNAMES = {0..N-1}
6
7  channel sits, eats, getsup : PHILNAMES
8  channel picks, puttdown : PHILNAMES.FORKNAMES
9
10 PHIL(i) = sits!i -> picks!i!i -> picks!i!((i+1)%N) ->
11 |     |     |     eats!i -> puttdown!i!((i+1)%N) -> puttdown!i!i -> getsup!i -> PHIL(i)
12
13 FORK(i) = picks!i!i -> puttdown!i!i -> FORK(i)
14 |     |     |     [] picks!((i-1)%N)!i -> puttdown!((i-1)%N)!i -> FORK(i)
15
16 PHILS = ||| i:PHILNAMES @ PHIL(i)
17 FORKS = ||| i:FORKNAMES @ FORK(i)
18
19 SYSTEM' = PHILS[|{|picks, puttdown|}|]FORKS
20
21 -- The potential for deadlock is illustrated by
22
23 assert SYSTEM' :[deadlock free [F]]
24

```

⁴Code from [FDR Documentation](#)



CSP: state space



$N = 2$

- States: 36
- Transitions: 65

$N = 3$

- States: 214
- Transitions: 565

$N = 4$

- States: 1,296
- Transitions: 4,569

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Timed automata

Focus: timed systems

Application: numerous applications
(<https://uppaal.org/casestudies/>)

Tools

- UPPAAL: <https://uppaal.org/>



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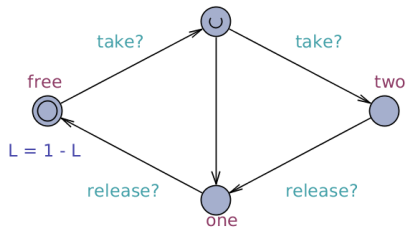
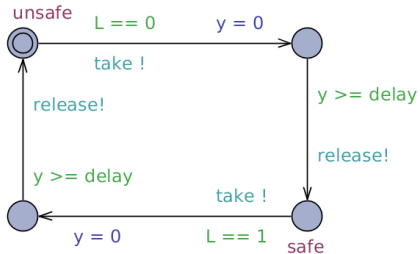
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Timed automata: example

Bridge crossing puzzle (Vikings version)



Timed automata: specification⁵



⁵Source: UPPAAL demos



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SMT-LIB v2

“*SMT-LIB* is an international initiative aimed at facilitating research and development in *Satisfiability Modulo Theories* (SMT).”

Focus: determining whether a mathematical formula is satisfiable
(https://en.wikipedia.org/wiki/Satisfiability_modulo_theories)

Application: numerous applications

Tools: numerous tools

- Z3:

www.microsoft.com/en-us/research/project/z3-3/



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SMT-LIB v2: example

```

3x + 2y = 6      (declare-const x Real)
2x - 3y + 5     (declare-const y Real)
                 (assert (= (+ (* 3 x) (* 2 y)) 6))
                 (assert (= (- (* 2 x) (* 3 y)) 5))
                 (check-sat)
                 (get-model)
                 (exit)

```

Output:

```

sat
(
  (define-fun y () Real (- (/ 3.0 13.0)))
  (define-fun x () Real (/ 28.0 13.0))
)

```

SMT-LIB v2: example

```
(p => q) ^ (q => p)  (declare-const p Bool)
  ¬ (p = q)         (declare-const q Bool)
                   (assert (and (=> p q) (=> q p)))
                   (assert (not (= p q)))
                   (check-sat)
                   (exit)
```

Output:

unsat



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Gallina

Focus: general purpose specification language

Application: numerous applications

(DeepSpec: <https://deepspec.org/main>)

Tools

- Coq proof assistant: <https://coq.inria.fr/>



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Gallina: example

Implementing dictionary ADT using binary search trees

```
1 clear      : dictionary -> dictionary
2 is_empty  : dictionary -> bool
3 size      : dictionary -> nat
4 insert    : key -> val -> dictionary -> dictionary
5 find      : key -> dictionary -> option val
```



Gallina: specification, implementation and proofs

```

21 Module Type Dictionary.
22
23 (* [key]: the type of keys *)
24 Definition key := nat.
25 (* [val]: the type of values *)
26 Definition val := string.
27 (* [dictionary]: the type of dictionaries *)
28 Parameter dictionary : Type.
29 (* [empty_dictionry]: definition of an empty dictionary *)
30 Parameter empty_dictionary : dictionary.
31
32 (**
33 Operations supported: [clear], [is_empty],
34 [size], [insert], [find]
35 *)
36
37 Parameter clear : dictionary -> dictionary.
38 Parameter is_empty : dictionary -> bool.
39 Parameter size : dictionary -> nat.
40 Parameter insert : key -> val -> dictionary -> dictionary.
41 Parameter find : key -> dictionary -> option val.
--

```



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Conclusion

Mathematical approaches to support the **rigorous** specification, design and verification of computer systems

In this talk:

- **B language**: Atelier B, ProB, BMotionWeb
- **CSP**: FDR4
- **Timed automata**: UPPAAL
- **SMT-LIB v2**: Z3
- **Gallina**: Coq



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Recommended reading: papers

Formal methods in dependable systems engineering: a survey of professionals from Europe and North America

Mario Gleirscher, Diego Marmsoler.

Empirical Software Engineering 25, pages 4473–4546. Springer 2020.

The 2020 Expert Survey on Formal Methods

Hubert Garavel, Maurice H. ter Beek, Jaco van de Pol.

Proc. Formal Methods for Industrial Critical Systems (FMICS 2020).

LNCS 12327, Springer 2020

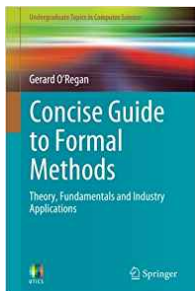


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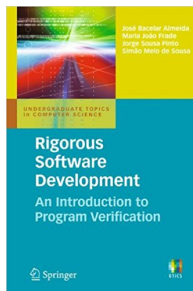
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Recommended reading: books



Gerard O'Regan

Concise Guide to Formal Methods
Springer, 2017.



José Almeida et al.

Rigorous Software Development
Springer, 2011.

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